## Modeling and Mitigating the Coremelt Attack

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#### Introduction

The Coremelt attack on a TCP network with the "dumbbell" topology



- Contribution
  - A dynamical system model for analysis
  - A limited number of subverted machines (bots): a modified TCP algorithm
  - A flow-based mitigation method
  - Simulation results

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## Distributed denial of service (DDoS) attack

 Attempt to disrupt network service by sending superfluous traffics from a vast number of bots

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- Soaring number of Internet of Things (IoT)  $\implies$  Escalating DDoS threats
  - 21 billion IoT devices by 2020

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## Distributed denial of service (DDoS) attack

- Attempt to disrupt network service by sending superfluous traffics from a vast number of bots
- Soaring number of Internet of Things (IoT)  $\implies$  Escalating DDoS threats
  - 21 billion IoT devices by 2020
- One of world's largest DDoS attack to date [Ant+17]
  - 2016 on OVH (hosting service in France)
  - Mirai Botnet: 150,000 hacked IoT devices, 600,000 at peak
  - Attack flow rate: 1 Tbps

[Ant+17] M. Antonakakis, T. April, M. Bailey, M. Bernhard, E. Bursztein, J. Cochran, Z. Durumeric, J. A. Halderman, L. Invernizzi, M. Kallitsis, D. Kumar, C. Lever, Z. Ma, J. Mason, D. Menscher, C. Seaman, N. Sullivan, K. Thomas, and Y. Zhou, in 26th USENIX Secure Symp., 2017 = 2000

#### The Coremelt attack



- A link-flooding DDoS attack [SP11]
- Target: backbone link

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- Target: backbone link
- Distributed botnet
  - Available
    - Mirai Botnet: 150k bots, 600k at peak
    - Among M bots there are  ${\cal O}(M^2)$  connections
  - Affordable
    - Price per 1000 bots: 100-180 in U.S. or U.K., 20-60 in Europe, less than 10 elsewhere

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Yang et al. (UCSB, UW, WPI)

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    - Among M bots there are  ${\cal O}(M^2)$  connections
  - Affordable
    - Price per 1000 bots: 100-180 in U.S. or U.K., 20-60 in Europe, less than 10 elsewhere
- Low-intensity, legitimate-looking traffic
  - Able to evade conventional DDoS defenses

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## Transmission Control Protocol (TCP)

- A congestion control algorithm [Pos81]
  - One congestion window per round-trip time (RTT)
  - Detect congestion based on missing acknowledgements (ACKs)
  - Additive-increase/multiplicative-decrease (AIMD) feedback algorithm [CJ89]

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- TCP-NewReno [Hen+12]
  - Widely used in modern Internet
  - Better for bursts of packet drops

[Pos81] J. Postel, Information Sciences Institute, Tech. Rep., 1981

<sup>[</sup>CJ89] D.-M. Chiu and R. Jain, Comput. Networks ISDN Syst., 1989

#### Dynamical system model

- Analyze the impact and effectiveness of the Coremelt attack
- Establish flow composition and convergence via Lyapunov-based analysis
- Understand the relations between the number of bots, packet drop probability, and link usage ratio of users
- Develop a flow-based mitigation method

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#### **TCP-NewReno source**

- One congestion window  $w_k$  per RTT  $\tau_k$
- Average flow rate  $x_k = w_k / \tau_k$
- Congestion probability  $q_k \approx w_k p$  with packet drop probability p

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- AIMD algorithm for TCP-NewReno

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Dynamical system model:

$$\dot{x}_k = \frac{1}{\tau_k^2} \left( (1 - q_k) - \frac{w_k}{2} q_k \right)$$

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#### **TCP-NewReno source**

$$\dot{x}_k = \frac{1 - \tau_k x_k p}{\tau_k^2} - \frac{p x_k^2}{2}, \qquad k = 1, \dots, N$$

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#### **Bottleneck link**

• Aggregate rate 
$$y = \sum x_k$$

- Bandwidth C
- Drop the excess packets

$$p = \begin{cases} 1 - C/y, & \text{if } y > C; \\ 0, & \text{otherwise} \end{cases}$$

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## Attack with M bots following TCP-NewReno



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#### Theorem 1

- If M bots and N M users all follow TCP-NewReno, the dynamical system is globally asymptotically stable (GAS)
- Packet drop probability converge to  $p^*$  satisfying  $\sum_{k=1}^{N} \frac{1}{\tau_k} = \frac{\sqrt{1+2/p^*+1}}{2(1-p^*)} p^*C$

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#### **Proof**

 $\blacksquare$  Lyapunov function  $V(x-x^{\ast})$  such that

$$\dot{V}(x-x^*) \le -W(x-x^*) - (p-p^*)(y-y^*)$$

- $W(x x^*)$  is positive definite
- $\blacksquare$  Packet drop probability p is increasing in aggregate rate y

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#### Implication

 $\blacksquare$  For the same RTT  $\tau,$  the link usage ratio of users is 1-M/N

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#### Implication

- $\blacksquare$  For the same RTT  $\tau,$  the link usage ratio of users is 1-M/N
- A target value  $p^*$  can be achieved by enough bots so that  $N \geq \frac{\sqrt{1+2/p^*}+1}{2(1-p^*)}\,p^*\tau C$

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## Attack with M bots following a modified TCP



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## Attack with M bots following a modified $\mathsf{TCP}$



#### Modified TCP source

- Internal state  $\xi_j$  that follows the AIMD algorithm for TCP-NewReno
- Flow rate  $x_j = \lambda_j \xi_j$  with gain  $\lambda_j \ge 0$
- Drive the congestion probability to target value  $q_0$  by slowly adjusting  $\lambda_j$ :

$$\dot{\lambda}_j = \gamma_j \xi_j (q_0 - q_j)^+_{\lambda_j}$$

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## Attack with M bots following a modified $\mathsf{TCP}$



#### Theorem 2

- If N M users follow TCP-NewReno and M bots follow the modified TCP, the dynamical system is GAS
- Congestion probability converge to target value  $q_0$  for any M

## Attack with M bots following a modified $\mathsf{TCP}$



#### Theorem 2

- If N M users follow TCP-NewReno and M bots follow the modified TCP, the dynamical system is GAS
- Congestion probability converge to target value  $q_0$  for any M

#### <u>Proof</u>

 $\blacksquare$  Weak Lyapunov function  $V(x_u-x_u^*,\xi-\xi^*,\lambda-\lambda^*)$  such that

$$\dot{V}(x_u - x_u^*, \xi - \xi^*, \lambda - \lambda^*) \le -W(x_u - x_u^*, \xi - \xi^*) - (p - p^*)(y - y^*)$$

- $W(x_u x_u^*, \xi \xi^*)$  is positive definite, p is increasing in y
- LaSalle's invariance principle

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## Mitigation

Detection-based mitigation: source authentication, packets inspection

- Less effective against Coremelt:
  - Communication between bot pairs
  - Low-intensity, legitimate-looking traffic

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  - Monitor source flow rates and assign individual drop probability  $p_k$  so that the bandwidth C is evenly shared:  $p_k \sim 1 C/(Nx_k)$

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  - Advantages:
    - Guaranteed link usage ratio of users: 1-M/N
    - Does not require modifying source transmission protocols
  - Limitations:
    - Extra resources needed to monitor source flow rates
    - Users with smaller RTTs will also be penalized
    - No effect against attacks with bots following TCP-NewReno

## Simulation: without mitigation

- $\blacksquare$  Network of 2,000 users and 1,000 bots
- Link capacity of 1 million packets per RTT



- Attack with TCP-NewReno: low congestion probability; link usage ratio of users is 2/3
- Attack with modified TCP: target congestion probability; link usage ratio of users is low

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### Simulation: with mitigation

- $\blacksquare$  Network of  $2000~{\rm users}$  and  $1000~{\rm bots}$
- Link capacity of 10<sup>6</sup> packets per RTT



 Attack with modified TCP: target congestion probability; link usage ratio of users is high

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#### Conclusion

- Contribution
  - A dynamical system model for analyzing the Coremelt attack on a TCP network
  - A limited number of bots: a modified TCP algorithm
  - A flow-based mitigation method
  - Simulation results

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  - A limited number of bots: a modified TCP algorithm
  - A flow-based mitigation method
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- Future work
  - User Datagram Protocol (UDP) [Pos80]
  - The Crossfire attack [KLG13]

[Pos80] J. Postel, Information Sciences Institute, Tech. Rep., 1980

[KLG13] M. S. Kang, S. B. Lee, and V. D. Gligor, in 2013 IEEE Symp.⊡Secur⊕Priv., 2013( ≥) ≥ - ∞ <

#### Conclusion

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#### Acknowledgements





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